

DISCRETE MATH - CS275 MIDTERM 2 - SOLUTION

NOVEMBER 16TH 2004, 12:30-13:45

Each of the 4 exercises is worth 5 points.

No calculator is needed or allowed in this exam.

All answers must come with an explanation.

Advice:

- Read all the exercises before answering to any.
- Spend 10 to 15 minutes on each exercise, then return to questions you missed.
- For each exercise, read **each word** with the greatest care and **without hurrying**. If needed, read it **many times**, until the meaning of the questions is clear.
- Check for extra information on the blackboard.

Exercise 1. Let F , G and H be Boolean functions of degree four, defined by

$$\begin{aligned}F(w, x, y, z) &= w + x + y + z \\G(w, x, y, z) &= \bar{w}xyz + w\bar{x}yz + wx\bar{y}z + wxy\bar{z} \\H(w, x, y, z) &= wx + wy + wz + xy + xz + yz\end{aligned}$$

Recall that, for any Boolean functions A , B of degree 4, one writes $A \leq B$ if and only if $\forall w, x, y, z \in \{0, 1\}$, $A(w, x, y, z) \leq B(w, x, y, z)$. Also recall that, if $A \leq B$, then $A(w, x, y, z)B(w, x, y, z) = A(w, x, y, z)$ for all $w, x, y, z \in \{0, 1\}$.

- a) Prove or disprove each of the following: $F \leq G$, $F \leq H$, $G \leq F$, $G \leq H$, $H \leq F$ and $H \leq G$.

Solution:

- 1) $F \leq G$: False: $F(0, 0, 0, 1) = 1$, $G(0, 0, 0, 1) = 0$.
- 2) $F \leq H$: False: $F(0, 0, 0, 1) = 1$, $H(0, 0, 0, 1) = 0$
- 3) $G \leq F$: True. Seen from truth table.
- 4) $G \leq H$: True. Seen from truth table.
- 5) $H \leq F$: True. Seen from truth table.
- 6) $H \leq G$: False: $H(0, 0, 1, 1) = 1$, $G(0, 0, 1, 1) = 0$

- b) Design a circuit using inverters, AND and OR gates, that implements the Boolean function

$$L(w, x, y, z) = F(w, x, y, z)G(w, x, y, z)H(w, x, y, z).$$

You may, if you wish and if you provide an explanation, simplify this Boolean expression to obtain a simpler circuit producing the same output.

Solution:

Since $G \leq F$, one has $FG = G$. Since moreover $G \leq H$, one has $GH = G = (FG)H = FGH$. So one has $L = G$ and the circuit in Fig. implements L .

It seems everyone knows how to draw the circuit of a Boolean circuit (good!), so I do not include a figure here.

Exercise 2. A mailbox contains 1,000,000 e-mails, each one having from 1 word to 1,000 words. Altogether, the mails have a total of 54,000,000 words. In the mails, 50,000 distinct words can be found. One of these words is "hello."

Prove or disprove each of the statements below:

- a) One mail is at least 55 words long.

Solution: No: Proof by counter-example. Consider a mailbox of 1,000,000 mails of length 54, so that $50e3$ distinct words are present. The total number of words is $54e6$ and no mail has 55 or more words.

- b) At least a thousand of the mails contain the same number of words.

Solution: Yes: Proof using the pigeonhole principle. Let w_i , for $i \in \{1, \dots, 1000\}$ be the number of mails containing i words. One has $w_1 + w_2 + \dots + w_{1000} = 1,000,000$, so that, by the pigeonhole principle, $w_i \geq \lceil 1,000,000/1,000 \rceil = 1000$ for some i .

- c) One word has a total number of occurrences, in all mails, of 1000 or more.

Solution: Yes: use the pigeonhole principle. Let y_n ($n \in \{1, \dots, 50,000\}$) be the total number of occurrences of word n . Since $y_1 + \dots + y_{50,000} = 54e6$, the pigeonhole principle says that the number of occurrences y_n of one word at least is $\lceil 54e6/5e5 \rceil = 1,080$ or greater.

- d) Two mails, at least, contain both the same number of words and the same number of occurrences of the word "hello."

Solution: Yes. Proof using the pigeonhole principle. Let $p = (x, z)$ be, a pair of numbers, representing respectively the number of words and the number of occurrences of "hello" in a message. One has $0 \leq z \leq x \leq 1000$ and also $0 < x$. The possible pairs are thus $p_1 = (1, 0)$, $p_2 = (1, 1)$, $p_3 = (2, 0)$, \dots , $p_{500,499} = (1000, 1000)$. Let u_i be the number of mails for which the number of words x and occurrences of "hello" z are equal to p_i . Since the u_i sum up to the number of e-mails, i.e. $u_1 + \dots + u_{500,499} = 1e6$, by the pigeonhole principle, one of the u_i must be greater or equal to $\lceil 1e6/500,499 \rceil = 2$.

Exercise 3. Solve the following recurrence relations:

- a) $x_n = -\frac{3}{4}x_{n-1} - \frac{1}{8}x_{n-2}$, with initial conditions $x_0 = x_1 = 1$.

Solution: The characteristic polynomial $r^2 + \frac{3}{4}r + \frac{1}{8}$ has distinct roots $r_1 = -\frac{1}{4}$ and $r_2 = -\frac{1}{2}$. From the initial conditions, one gets

$$x_n = 6 \left(-\frac{1}{4}\right)^n - 5 \left(-\frac{1}{2}\right)^n$$

- b) $y_n = 2y_{n-1} - \frac{8}{9}y_{n-2}$, with initial conditions $y_0 = y_1 = 1$.

Solution: The characteristic polynomial $r^2 - 2r + \frac{8}{9}$ has distinct roots $r_1 = \frac{4}{3}$ and $r_2 = \frac{2}{3}$. From the initial conditions, one gets

$$y_n = \frac{1}{2} \left(\frac{4}{3}\right)^n + \frac{1}{2} \left(\frac{2}{3}\right)^n.$$

One may check that $y_2 = \frac{10}{9}$, $y_3 = \frac{12}{9} = \frac{4}{3}$.

Exercise 4. Let $A = \{0, 1\}^5$ be the set of binary sequences (strings) of zeros and ones of length 5. We will write $X = (x_1, x_2, \dots, x_5)$ an element of A , where $x_i \in \{0, 1\}$ for all $i \in \{1, \dots, 5\}$. Also, e.g. $(1, 0, 1, 0, 0)$ will be written 10100.

- a) Let f be the function defined by

$$f : \begin{cases} A & \longrightarrow \{0 \dots 5\} \\ X & \longrightarrow x_1 + \dots + x_5, \end{cases}$$

and let R be the relation defined on A by

$$R(X, Y) \equiv f(X) = f(Y).$$

- 1) Prove or disprove that f is onto.

Solution: Yes: For any i in the co-domain $\{0, \dots, 5\}$ of f , on has $f(\underbrace{1\dots 1}_i \underbrace{0\dots 0}_{5-i}) = i$, so

that f is onto.

- 2) Prove or disprove that f is one-to-one.

Solution: No: $f(1, 0, \dots, 0) = f(0, 1, 0, 0, 0) = 1$.

- 3) Prove that R is an equivalence relation.

Solution:

- R is (trivially) reflexive, since $\forall X, f(X) = f(X)$.
- R is (trivially) symmetric since for all X, Y , one has $f(X) = f(Y) \implies f(Y) = f(X)$.
- R is (trivially) transitive since for all X, Y, Z , one has $f(X) = f(Y) \wedge f(Y) = f(Z) \implies f(X) = f(Z)$.

- 4) Write the equivalence classes of 01010 and of 10000.

Solution:

- $\bar{X}_1 = \{X \in A \mid f(X) = 2\} = \{11000, 10100, 10010, 10001, 01100, 01010, 01001, 00110, 00101, 00011\}$.
- $\bar{X}_2 = \{X \in A \mid f(X) = 1\} = \{10000, 01000, 00100, 00010, 00001\}$.

- 5) Write the cardinal of the equivalence class of X as a function of $f(X)$.

Solution: $|\bar{X}| = C(5, f(X))$.

- b) Let S be the relation defined on A by

$$S(X, Y) \equiv \forall i \in \{1, \dots, 5\}, x_i \leq y_i.$$

- 1) Prove that S is a partial order relation.

Solution:

- S is (trivially) reflexive, since $\forall X, \forall i \in \{1, \dots, 5\} x_i \leq x_i$.
- R is (trivially) antisymmetric since, for all X, Y , if $\forall i \in \{1, \dots, 5\} x_i \leq y_i$ and $\forall i \in \{1, \dots, 5\} y_i \leq x_i$, then $\forall i, x_i = y_i$.
- R is (trivially) transitive since for all X, Y, Z , if $\forall i \in \{1, \dots, 5\} x_i \leq y_i$ and $\forall i \in \{1, \dots, 5\} y_i \leq z_i$, then $\forall i, x_i \leq z_i$.

- 2) Prove or disprove that S is a total order.

Solution: Let $Y_1 = (1, 0, \dots, 0) \in A$ and $Y_2 = (0, 1, 0, \dots, 0) \in A$. Since $\neg S(Y_1, Y_2)$ and $\neg S(Y_2, Y_1)$, i.e. Y_1 and Y_2 are not comparable, S is not a total order.

- 3) Write the set of elements $Y \in A$ such that $S(Y, 10100)$.

Solution: $\{00000, 10000, 00100, 10100\}$.

- 4) Write the set of upper and lower bounds of the set $\{10100, 01010\} \subseteq A$.

Solution: Upper bounds: $\{11110, 11111\}$. There is a single lower bound: 00000.

- 5) Write the sets of maximal and of minimal elements of A and, if they exist, the greatest and least element of A .

Solution: It is clear that $(\underbrace{1, \dots, 1}_n)$ is greater than all other elements and is thus the greatest element and the only maximal element; and that $(\underbrace{0, \dots, 0}_n)$ is less than all other elements and is thus the least element and the only minimal element.